Finite Automata and Formal Languages

Final Examination

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CSIE210030, National Chi Nan University

Problem 1 Prove that the class of context-free languages is not closed under the complement. (Hint: Show that $A = \{a^m b^n c^n | m, n \ge 0\}$ and $B = \{a^n b^n c^m | m, n \ge 0\}$ are context free but their intersection is not.)

Problem 2 State the Pumping Lemma for context-free languages. Show that $L = \{a^n b^m | m \ge n^2\}$ is not a context-free language.

Problem 3 Convert the context-free grammar

$$S \rightarrow Sa \mid SbSa \mid \epsilon$$

into an equivalent (nondeterministic) pushdown automaton.

Problem 4 Let \mathcal{B} be the set of all infinite sequences over $\{0, 1\}$. Show that B is uncountable, using a proof by the diagonalization method.

Problem 5 Let $ONE_{DFA} = \{\langle A \rangle \mid A \text{ is a DFA over } \Sigma \text{ that recognizes } 1^* \}$ where the alphabet $\Sigma = \{0, 1\}$. Show that ONE_{DFA} is decidable.

Problem 6 Let $E_{CFG} = \{\langle G \rangle \mid G \text{ is a CFG and } L(G) = \emptyset\}$. Show that E_{CFG} is decidable.

Problem 7 Let $A_{TM} = \{ \langle M, w \rangle | M \text{ is a Turing machine and } M \text{ accepts } w \}$. Show that A_{TM} is undecidable.

Problem 8 Let $M_{\epsilon} = \{\langle M \rangle | M \text{ is a Turing machine and } M \text{ halts on } \epsilon\}$. Explain why the following algorithm is not a description of a legitimate Turing machine that decides M_{ϵ} . $D_{\epsilon} = \text{``On input } M,$

- 1. IF M halts on ϵ THEN Accept;
- 2. ELSE Reject."

